

Logic and Foundations II

Part 8. Second order arithmetic and non-standard methods

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Logic and Foundations II

- Part 5. Models of first-order arithmetic (continued) (5 lectures)
- Part 6. Real-closed ordered fields: completeness and decidability (4 lectures)
- Part 7. Real analysis and reverse mathematics (8.5 lectures)
- **Part 8. Second order arithmetic and non-standard methods** (6.5 lectures)

Part 8. Schedule

- May 21, (0) Introduction to forcing
- May 23, (1) Harrington's conservation result on WKL_0
- May 28, (2) H.Friedman's conservation result on WKL_0
- May 30, (3) Friedman's result (continued) and a self-embedding theorem I
- **June 04, (4) A self-embedding theorem II**
- June 06, (5)
- June 11, (6)

§8.3. A self-embedding theorem of WKL_0

In this section, we introduce a self-embedding theorem of WKL_0 , by which we can devise methods of nonstandard analysis in WKL_0 .

Gödel stated in 1973 that "nonstandard analysis is the future of analysis." However, Henson and Keisler have shown in 1986 that nonstandard arguments in n -th order arithmetic require $(n + 1)$ -th order arithmetic. Therefore, conducting complete nonstandard analysis for second-order arithmetic Z_2 is impossible within the framework of second-order arithmetic alone. Nevertheless, as demonstrated in my paper¹, certain amount of nonstandard analysis can still be developed within WKL_0 .

The main tool of our nonstandard method is a self-embedding theorem of WKL_0 (Theorem 3.9), which extends Friedman's self-embedding theorem (§5.3) to WKL_0 . This section primarily discusses the proof of this theorem.

¹K. Tanaka, The self-embedding theorem of WKL_0 and a non-standard method, *Annals of Pure and Applied Logic* 84 (1997), pp.41–49.

Theorem 3.1 (Self-Embedding Theorem)

Let $\mathfrak{M} = (M, S)$ be a countable model of WKL_0 with $M \neq \omega$. Then, there exists a proper initial segment I of M such that $\mathfrak{M}[I = (I, S[I)$ is isomorphic to \mathfrak{M} . Here, $S[I = \{X \cap I \mid X \in S\}$.

Before proving this theorem, we need some preparations. We first prove the following lemma, which will be frequently used later.

Lemma 3.2 (Compactness in WKL_0)

(1) For any Π_1^0 formula $\varphi(X)$, there exists a Π_1^0 formula $\hat{\varphi}$ such that WKL_0 proves:

$$\hat{\varphi} \leftrightarrow \exists X \varphi(X).$$

(2) For any Π_1^0 formula $\varphi(k, X)$, WKL_0 proves:

$$\forall n \exists X \forall k < n \varphi(k, X) \rightarrow \exists X \forall k \varphi(k, X).$$

From now on, we adopt the notation $[T]$ for the set of all infinite paths of a tree T . Do not confuse it with $[p]$, which represents a basic open set in the order topology.

Proof. (1) We identify a set X with its characteristic function, which is also represented as an infinite binary sequence. Then, a Π_1^0 formula $\varphi(X)$ can be expressed as $\forall x \theta(X \upharpoonright x)$, where θ is Σ_0^0 and $X \upharpoonright x$ is a code for a finite binary sequence. We set $T = \{t \mid \forall s \subseteq t \theta(s)\}$. Then T is a tree, and $X \in [T]$ iff $\varphi(X)$ holds. Thus, $\exists X \varphi(X)$ is equivalent to $[T] \neq \emptyset$, which is expressed as a Π_1^0 formula “ T is infinite ($\forall n \exists t \in \{0, 1\}^n t \in T$)”.

(2) Express a Π_1^0 formula $\varphi(k, X)$ as $\forall x \theta(k, X \upharpoonright x)$ (where θ is Σ_0^0), and define a tree $T = \{t \mid \forall k \leq \text{length}(t) \forall x \leq \text{length}(t) \theta(k, t \upharpoonright x)\}$. Here, $\text{length}(t)$ denotes the length of the finite binary sequence t . If $\forall n \exists X \forall k < n \varphi(k, X)$ holds, then $\forall n \exists X \forall k < n \forall x < n \theta(k, X \upharpoonright x)$, so $t = X \upharpoonright n \in T$ for all n , thus T is infinite. Hence, in WKL_0 , there exists an infinite path $X \in [T]$ satisfying $\forall k \varphi(k, X)$. \square

Here is another demonstration for (2). If we express $\varphi(k, X)$ as $X \in [T_k]$, then $\exists X \forall k < n \varphi(k, X)$ can be expressed as $\bigcap_{k < n} [T_k] \neq \emptyset$. Since this is true for any n , we have $\bigcap_{k < \infty} [T_k] \neq \emptyset$ by the compactness of the Cantor space since $[T_k]$'s are closed sets.

Both (1) and (2) are referred to as “compactness (of binary trees) in WKL_0 ”.

We define $G\text{-}\Sigma_1^0$ **formulas** or simply G **formulas** by generalizing Σ_1^0 formulas as follows. The G formulas are obtained from Σ_1^0 formulas by using \wedge, \vee , bounded universal quantifier $\forall x < y$ and unbounded existential quantifier $\exists x$, and set quantifiers $\forall X, \exists X$.

In WKL_0 , we can prove that a G formula is equivalent to a Σ_1^0 formula.

(Proof)

- The closure condition under $\forall x < y$ is nothing but the collection principle $\text{B}\Sigma_1^0$ derivable from Σ_1^0 induction.
- The closure condition under $\forall X$ can be obtained from Lemma 3.2(1) by taking the negation on both sides.
- The closure condition under $\exists X$ can be demonstrated by noting that $\exists X \exists x \theta(x, X \upharpoonright x)$ (where θ is Σ_0^0) can be rewritten as $\exists t \exists x \theta(x, t)$.
- The other closure conditions are almost obvious.

Now, we redefine the G -formulas explicitly in RCA_0 in the next slide.

Definition 3.3 (G -formulas)

A sequence $G_0 \subset G_1 \subset G_2 \subset \dots$ of sets of $\mathcal{L}_{\text{OR}}^2$ -formulas is defined inductively modulo 4 as follows: for each $e \in \mathbb{N}$,

$$\begin{aligned} G_0 &= \{\text{finite disjunctions } (\vee) \text{ of atomic formulas or their negations}\}, \\ G_{4e+1} &= \{\exists x \phi \mid \phi \text{ is a finite conjunction } (\wedge) \text{ of } G_{4e} \text{ formulas}\} \cup G_{4e}, \\ G_{4e+2} &= \{\forall x < y \phi \mid \phi \text{ is a finite disjunction } (\vee) \text{ of } G_{4e+1} \text{ formulas}\} \cup G_{4e+1}, \\ G_{4e+3} &= \{\exists X \phi \mid \phi \text{ is a finite conjunction } (\wedge) \text{ of } G_{4e+2} \text{ formulas}\} \cup G_{4e+2}, \\ G_{4e+4} &= \{\forall X \phi \mid \phi \text{ is a finite disjunction } (\vee) \text{ of } G_{4e+3} \text{ formulas}\} \cup G_{4e+3}. \end{aligned}$$

Finally, we set $\mathbf{G} = \bigcup_{e \in \mathbb{N}} G_e$. The formulas in \mathbf{G} are called **G formulas**.

By Lemma 5.5.3, there is no formula that defines the truth values of all formulas. But, Lemma 5.3.4 shows that if we restrict the formulas to a class like Σ_n , then there exists a formula Sat_{Σ_n} to define the truth values of formulas in the class. This is also the case for Σ_n^0 in second order arithmetic. In the following, we will define Sat for G formulas.

From now on, a structure $\mathfrak{M} = (M, S)$ is denoted by V . Then, for each $p \in M$, set $M_p = \{a \in M \mid \mathfrak{M} \models a < p\}$, $S_p = \{X \cap M_p \mid X \in S\}$ and denote $V_p = (M_p, S_p)$.

Since M_p may not be closed under operations such as addition, V_p may not be a substructure of V . However, just by restricting the ranges of variables to these sets, the **satisfaction predicate** $\text{Sat}^p(z, \xi)$ for V_p can be naturally defined within $V = (M, S)$. Here, z represents the code of a formula φ , and ξ is a finite function that assigns elements of $M_p \cup S_p$ to free variables appearing in φ . Thus, supposing that a formula $\varphi(\vec{x}, \vec{X})$ has no free variables other than \vec{x}, \vec{X} , and $\xi(\vec{x}) = \vec{a}, \xi(\vec{X}) = \vec{U}$, we have in V ,

$$\text{Sat}^p(\ulcorner \varphi \urcorner, \xi) \equiv \varphi(\vec{a}, \vec{U})^{V_p}, \text{ roughly } V_p \models \varphi(\vec{a}, \vec{U}).$$

Here, in $\varphi(\vec{a}, \vec{U})^{V_p}$, quantification over numbers is bounded by p , and quantification over sets is also considered as ranging binary sequences of length p , which can be coded by numbers $< 2^p$. Thus, $\text{Sat}^p(z, \xi)$ can be defined as a Δ_1^0 formula in V (cf. Lemma 5.3.4).

We also remark that a variable z in $\text{Sat}^p(z, \xi)$ can potentially express a non-standard number. In V , it can be easily verified that Sat^p satisfies Tarski's truth definition clauses for all standard formulas (cf. Theorem IV.2.26 in [P. Hájek and P. Pudlák, *Metamathematics of First-order Arithmetic*, Springer, 1993.]).

Next, we define the **satisfaction relation for G formulas** as follows:

Definition 3.4

For each $z \in G$, define the satisfaction relation $\text{Sat}(z, \xi)$ as follows:

$$\text{Sat}(z, \xi) \leftrightarrow \exists p \text{Sat}^p(z, \xi \upharpoonright V_p).$$

Here, $\xi \upharpoonright V_p$ is the assignment obtained by restricting the values of ξ to V_p .

For simplicity, we abbreviate $\text{Sat}^p(z, \xi \upharpoonright V_p)$ as $\text{Sat}^p(z, \xi)$. It is provable in RCA_0 that for the code z of a Σ_1^0 formula, if $\text{Sat}^p(z, \xi)$ holds, then $\text{Sat}^{p'}(z, \xi)$ also holds for any $p' \geq p$.

Moreover, we will show in WKL_0 that it also the case for the codes z of G .

In the following, we identify a formula with its code.

Lemma 3.5

In a model V of WKL_0 , $\text{Sat}(z, \xi)$ satisfies Tarski's truth definition clauses for G formulas.

Proof. We prove the statement by induction on the complexity of the formula z .

If z is an atomic formula or its negation, $\text{Sat}(z, \xi) \Leftrightarrow \exists p \text{Sat}^p(z, \xi) \Leftrightarrow \exists p z(\xi)^{V_p} \Leftrightarrow z(\xi)$.

If $z = \bigvee_{i < n} z_i$ (where each z_i is a G formula),

$$\begin{aligned} \text{Sat} \left(\bigvee_{i < n} z_i, \xi \right) &\Leftrightarrow \exists p \text{Sat}^p \left(\bigvee_{i < n} z_i, \xi \right) \Leftrightarrow \exists p \bigvee_{i < n} \text{Sat}^p(z_i, \xi) \\ &\Leftrightarrow \bigvee_{i < n} \exists p \text{Sat}^p(z_i, \xi) \Leftrightarrow \bigvee_{i < n} \text{Sat}(z_i, \xi). \end{aligned}$$

If z is $\exists x z'$ or $\exists X z'$ (where z' is a G formula), the proofs are analogous.

When $z = \bigwedge_{i < n} z_i$ (where each z_i is a G formula),

$$\begin{aligned} \text{Sat} \left(\bigwedge_{i < n} z_i, \xi \right) &\Leftrightarrow \exists p \text{Sat}^p \left(\bigwedge_{i < n} z_i, \xi \right) \Leftrightarrow \exists p \bigwedge_{i < n} \text{Sat}^p (z_i, \xi) \\ &\Leftrightarrow \bigwedge_{i < n} \exists p \text{Sat}^p (z_i, \xi) \quad (\Leftarrow \text{by } \Sigma_1^0 \text{ collection principle}) \\ &\Leftrightarrow \bigwedge_{i < n} \text{Sat} (z_i, \xi). \end{aligned}$$

If z is $\forall x < y z'$ (where z' is a G formula), the proof is analogous.

If $z = \forall X z'$ (where z' is a G formula),

$$\begin{aligned} \text{Sat} (\forall X z', \xi) &\Leftrightarrow \exists p \text{Sat}^p (\forall X z', \xi) \Leftrightarrow \exists p \forall U \text{Sat}^p (z', \xi \cup \{(X, U)\}) \\ &\Leftrightarrow \forall U \exists p \text{Sat}^p (z', \xi \cup \{(X, U)\}) \quad (\Leftarrow \text{by compactness (Lemma 3.2(2))}) \\ &\Leftrightarrow \forall U \text{Sat} (z', \xi \cup \{(X, U)\}), \end{aligned}$$

where $\xi \cup \{(X, U)\}$ is an extension of ξ with X assigned to U . □

Lemma 3.6

In a model $V = (M, S)$ of WKL_0 , we fix any $e \in M$ and an M -finite assignment map ξ . Then, there exists a $p \in M$ such that for all G_e formulas z whose free variables all belong to the domain of ξ , then $\text{Sat}(z, \xi) \Leftrightarrow \text{Sat}^p(z, \xi)$ holds.

Proof. Since the domain of the assignment map ξ is M -finite, the set of G_e formulas whose free variables are in the domain of ξ is essentially M -finite (disregarding repetitions of the same formulas within a disjunction or conjunction). This fact can be demonstrated by Σ_1^0 induction on e .

Therefore, for M -finitely many G_e formulas z , if $\text{Sat}(z, \xi)$ holds, let p_z be p such that $\text{Sat}^p(z, \xi)$, or otherwise let $p_z = 0$. Then, if we put $q = \max\{p_z\}$,² then we have $\text{Sat}(z, \xi) \Leftrightarrow \text{Sat}^q(z, \xi)$. □

²Strictly speaking, strong Σ_1^0 collection principle ($S\Sigma_1$) is used here. (Refer to Problem 1 following Lemma 1.8 in Chapter 7.)

Definition 3.7 (reflection)

In a model V of WKL_0 , for any e, p , and for two assignment maps ξ, ξ' with the same domain, the relation $\text{Ref}_e^p(\xi, \xi')$ is defined as follows:

$$\text{Sat}(z, \xi) \Rightarrow \text{Sat}^p(z, \xi'), \text{ for each } G_e \text{ formula } z \text{ with free variables in the domain of } \xi.$$

Lemma 3.8

In a model V of WKL_0 , supposing $\text{Ref}_e^p(\xi, \xi')$ with M -finite ξ, ξ' , the following holds:

- (1) If $e = 4d + 1$, $\forall a \exists a' < p \text{Ref}_{e-1}^p(\xi \cup \{(y, a)\}, \xi' \cup \{(y, a')\})$, where y is a variable not in the domain of ξ .
- (2) If $e = 4d + 2$, for each numerical variable x belonging to ξ ,
 $\forall a' < \xi'(x) \exists a < \xi(x) \text{Ref}_{e-1}^p(\xi \cup \{(y, a)\}, \xi' \cup \{(y, a')\})$, with y not in ξ .
- (3) If $e = 4d + 3$, $\forall U \exists U' \text{Ref}_{e-1}^p(\xi \cup \{(Y, U)\}, \xi' \cup \{(Y, U')\})$, where Y is a variable not belonging to the domain of ξ .
- (4) If $e = 4d + 4$, $\forall U' \exists U \text{Ref}_{e-1}^p(\xi \cup \{(Y, U)\}, \xi' \cup \{(Y, U')\})$, with Y not in ξ .

Proof Let $V = (M, S)$ be a model of WKL₀, and let ξ, ξ' be M -finite assignments with the same domain such that $\text{Ref}_e^p(\xi, \xi')$ is satisfied.

- (1) For $e = 4d + 1$. Show $\forall a \exists a' < p \text{Ref}_{e-1}^p(\xi \cup \{(y, a)\}, \xi' \cup \{(y, a')\})$.

Fix any $a \in M$. Let Z be the set of all codes of G_{e-1} formulas z satisfying $\text{Sat}(z, \xi \cup \{(y, a)\})$ and in a non-redundant form (i.e., no same formula is repeated in disjunctions or conjunctions), whose free variables are either y or belong to the domain of ξ . According to the argument in the proof of Lemma 3.6, this set Z is M -finite within V . Thus, by (bounded Σ_1^0 -CA) (Lemma 7.1.8), Z exists.

Now, consider a G_e -formula $z' = \exists y \bigwedge_{z \in Z} z$. Since $\text{Sat}(z, \xi \cup \{(y, a)\})$ for each $z \in Z$, it follows from Lemma 3.5 that $\text{Sat}(\bigwedge_{z \in Z} z, \xi \cup \{(y, a)\})$ and so $\text{Sat}(z', \xi)$.

Therefore, by the hypothesis, $\text{Sat}^p(z', \xi')$ holds. Thus, there exists $a' < p$ such that $\text{Sat}^p(z, \xi' \cup \{(y, a')\})$ holds for each $z \in Z$, fulfilling the requirement.

- (2) For $e = 4d + 2$. Show $\forall a' < \xi'(x) \exists a < \xi(x) \text{Ref}_{e-1}^p(\xi \cup \{(y, a)\}, \xi' \cup \{(y, a')\})$.

Fix any $a' < \xi'(x)$. To prove by contradiction, assume that for any $a < \xi(x)$ there exists a G_{e-1} formula z such that $\text{Sat}(z, \xi \cup \{(y, a)\})$ and $\neg \text{Sat}^p(z, \xi' \cup \{(y, a')\})$. Let Z be the set of all $z \in G_{e-1}$ satisfying $\neg \text{Sat}^p(z, \xi' \cup \{(y, a')\})$ and in a non-redundant form, whose free variables are either y or belong to the domain of ξ .

- (2) (continued) Like in case (1), Z exists by (bounded Σ_1^0 -CA). Consider a G_e formula $z' = \forall y < x \bigvee_{z \in Z} z$. By the other assumption, for each $a < \xi(x)$, there exists $z \in Z$ such that $\text{Sat}(z, \xi \cup \{(y, a)\})$, so $\text{Sat}(z', \xi)$ holds.

Therefore, by the hypothesis, $\text{Sat}^p(z', \xi')$ holds. Thus for each $a' < \xi'(x)$, there exists $z \in Z$ such that $\text{Sat}^p(z, \xi' \cup \{(y, a')\})$, which contradicts the definition of Z .

- (3) For $e = 4d + 3$. $\forall U \exists U' \text{Ref}_{e-1}^p(\xi \cup \{(Y, U)\}, \xi' \cup \{(Y, U')\})$ can be shown like (1).
- (4) For $e = 4d + 4$. Show $\forall U' \exists U \text{Ref}_{e-1}^p(\xi \cup \{(Y, U)\}, \xi' \cup \{(Y, U')\})$.
Fix any U' . Let Z be the set of $z \in G_{e-1}$ satisfying $\neg \text{Sat}^p(z, \xi' \cup \{(Y, U')\})$ and in a non-redundant form, whose free variables are either y or belong to the domain of ξ . Consider a G_e formula $z' = \forall Y \bigvee_{z \in Z} z$. By contradiction, assume for each U , there exists $z \in Z$ such that $\text{Sat}(z, \xi \cup \{(Y, U)\})$. Thus, $\text{Sat}(z', \xi)$ holds, and by the hypothesis, $\text{Sat}^p(z', \xi')$ holds, which contradicts the definition of Z .

Thus, the lemma is proved. □

Theorem 3.9 (Self-Embedding Theorem)

Let $\mathfrak{M} = (M, S)$ be a countable model of WKL₀ with $M \neq \omega$. Then, there exists a proper initial segment I of M such that $\mathfrak{M} \upharpoonright I = (I, S \upharpoonright I)$ is isomorphic to \mathfrak{M} .

Proof Let $V = (M, S)$ be a countable nonstandard model of WKL₀, and fix $q \in M$. Since V_q is M -finite within V , we can also make an M -finite mapping ξ_0 that assigns each number and set in V_q to distinct variables.

Now, take any nonstandard number $e \in M$. By Lemma 3.6, for any G_e -formula z whose free variables belong to the domain of ξ_0 , there exists p such that $\text{Sat}(z, \xi_0) \Leftrightarrow \text{Sat}^p(z, \xi_0)$ holds.

In the following, by repeatedly using Lemma 3.8 (the back-and-forth method), we construct two ω -sequences of assignment mappings $\xi_0 \subseteq \xi_1 \subseteq \dots \subseteq \xi_k \subseteq \dots$ and $\xi'_0 (= \xi_0) \subseteq \xi'_1 \subseteq \dots \subseteq \xi'_k \subseteq \dots$ ($k \in \omega$), where $\text{Ref}_{e-k}^p(\xi_k, \xi'_k)$ holds for all $k \in \omega$, and $\bigcup_k \text{range}(\xi_k) = V$ and $\bigcup_k \text{range}(\xi'_k)$ forms the desired initial segment of the model V .

To begin with, we enumerate the elements of V as $M = \{a_i \mid i \in \omega\}$, $S = \{U_i \mid i \in \omega\}$. We inductively construct ξ_k, ξ'_k with the same domain ($k \in \omega$) by cases:

- (i) For $e - k = 4d + 1$. Let a be the element a_i in $M - \text{range}(\xi_k)$ with the smallest index i , and let $a' < p$ be obtained by Lemma 3.8(1). Then, let y be a new numerical variable not in the domain of ξ_k , and set $\xi_{k+1} = \xi_k \cup \{(y, a)\}$, $\xi'_{k+1} = \xi_k \cup \{(y, a')\}$.

- (ii) For $e - k = 4d + 2$. Let $\xi'_k(x_0)$ be the largest in the order in M among all $\xi'_k(x)$'s. Then, let a' be the element a_i in $M - \text{range}(\xi'_k)$ and satisfying $a_i < \xi'_k(x_0)$ with the smallest index i , and let $a < \xi(x_0)$ be obtained by Lemma 3.8(2). Then, let y be a new numerical variable, and set $\xi_{k+1} = \xi_k \cup \{(y, a)\}$, $\xi'_{k+1} = \xi'_k \cup \{(y, a')\}$.
- (iii) For $e - k = 4d + 3$. Let U be $U_i \in S$ with the smallest index i , that is different from any set in $\text{range}(\xi_k)$ with regards to the numbers in $\text{range}(\xi_k)$. Also, let U' be obtained by Lemma 3.8(3). Then, let Y be a new set variable, and set $\xi_{k+1} = \xi_k \cup \{(Y, U)\}$, $\xi'_{k+1} = \xi'_k \cup \{(Y, U')\}$.
- (iv) For $e - k = 4d + 4$. Let U' be $U_i \in S$, with the smallest index i , that is different from any set in $\text{range}(\xi'_k)$ with regards to the numbers in $\text{range}(\xi'_k)$. Also, let U be obtained by Lemma 3.8(4). Then, let Y be a new set variable, and set $\xi_{k+1} = \xi_k \cup \{(Y, U)\}$, $\xi'_{k+1} = \xi'_k \cup \{(Y, U')\}$.

From the above construction, it is easy to see that $\text{Ref}_{e-k}^p(\xi_k, \xi'_k)$ holds for each $k \in \omega$.

From (i) and (iii), it is obvious that $\bigcup_k \text{range}(\xi_k) = (M, S)$. Also, from (ii), we can easily see that the set I consisting of a belonging to $\bigcup_k \text{range}(\xi'_k)$ forms an initial segment of M . Then, from (iv) it follows that $\bigcup_k \text{range}(\xi'_k) = (I, S \upharpoonright I)$.

Next, we prove by induction that both ξ_k, ξ'_k are injective for all $k \in \omega$. It is clear from the definition that $\xi_0 = \xi'_0$ is injective.

In (i), we first extend the injective mapping ξ_k to an injective ξ_{k+1} , and then extend the injective ξ'_k to a mapping ξ'_{k+1} that satisfies $\text{Ref}_{e-k-1}^p(\xi_{k+1}, \xi'_{k+1})$. The injectivity of ξ_{k+1} is clear from the construction. Since the injectivity is expressed by a G_2 formula, ξ'_{k+1} is also injective.

Similarly for (ii), (iii) and (iv).

Thus, $\bigcup_k \xi_k$ and $\bigcup_k \xi'_k$ are also injective.

Let $f = (\bigcup_k \xi'_k) \circ (\bigcup_k \xi_k)^{-1}$, which becomes a bijection from V to $V \upharpoonright I$. It is evident that f acts as the identity map on V_q .

Furthermore, since $\text{Ref}_0^p(\xi_k, \xi'_k)$ holds for each $k \in \omega$, it is clear that f is an isomorphism.

Thus, the proof of the theorem is complete. \square

Let's briefly describe how the Self-Embedding Theorem 3.9 can be applied to nonstandard analysis.

- According to Gödel's completeness theorem and compactness theorem,

$$WKL_0 \vdash \varphi \Leftrightarrow \text{for any non-}\omega \text{ model } \mathfrak{M} \text{ of } WKL_0, \mathfrak{M} \models \varphi.$$

- Since any infinite structure has an elementarily equivalent countable structure by the Löwenheim-Skolem Theorem,

$$WKL_0 \vdash \varphi \Leftrightarrow \text{for any countable non-}\omega \text{ model } \mathfrak{M} \text{ of } WKL_0, \mathfrak{M} \models \varphi.$$

- Choose a countable non- ω model $\mathfrak{M} = (M, S)$ of WKL_0 . Theorem 3.9 states that \mathfrak{M} has an initial segment isomorphic to itself. But by swapping their roles of \mathfrak{M} and an isomorphic initial segment, \mathfrak{M} is seen to have an isomorphic end-extension $^*\mathfrak{M} = (^*M, ^*S)$, which allows us to carry out some nonstandard analysis arguments.
- For example, in $\mathfrak{M} = (M, S)$, a real number a is indeed a set in S . Thus, a is an initial segment $^*a \upharpoonright M$ of some set $^*a \in ^*S$. Since *a may be taken bounded in $^*\mathfrak{M}$, it can be coded by an element of *M . Therefore, a real number in \mathfrak{M} can be treated like a rational number in $^*\mathfrak{M}$.

Application (The Maximum Principle)

$WKL_0 \vdash$ Any continuous function $f : [0, 1] \rightarrow [0, 1]$ has a maximum value.

Proof.

$$\mathfrak{M} = (M, S)$$

$$*\mathfrak{M} = (*M, *S)$$

$$\begin{array}{ccc}
 f : [0, 1] \cap \mathbb{Q} \rightarrow [0, 1] & \implies & *f : \{q_i\}_{i < a} \rightarrow 2^b \\
 \parallel & & (a, b \in *M - M, f = *f \cap M) \\
 \{q_i\}_{i \in M} & & 2^M
 \end{array}$$

$$\begin{array}{ccc}
 *m \cap M \text{ is sup } f & \longleftarrow & *m = \max\{ *f(q_i) \}_{i < a} \\
 & & \downarrow
 \end{array}$$

Other Applications

$WKL_0 \vdash$ The Cauchy-Peano Theorem (Tanaka, 1997)

$WKL_0 \vdash$ The existence of Haar measure for a compact group
(Tanaka-Yamazaki, 2000)

$WKL_0 \vdash$ The Jordan curve theorem (Sakamoto-Yokoyama, 2007)

Thank you for your attention!