

# *Topics in Applied Math:* Logic and Foundations of Mathematics

## Part 6. Reverse Mathematics

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November 28, 2025



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## Logic and Foundations

- **Part 1. Equational theory**
- **Part 2. First order theory**
- **Part 3. Basic Model theory**
- **Part 4. First order arithmetic and incompleteness theorems**
- **Part 5. Models of first-order arithmetic**
- **Part 6. Reverse mathematics**
- Part 7. Second order arithmetic and non-standard methods

## Part 6. Schedule

- Nov. 19, (1) Introduction and  $RCA_0$
- Nov. 21, (2) Real numbers in  $RCA_0$
- Nov. 26, (3) Completeness of the reals and  $ACA_0$
- **Nov. 28, (4) Continuous functions and  $WKL_0$**

**Reverse Mathematics:** Which axioms are needed to prove a theorem?

The Reverse Mathematics Phenomenon

*Many theorems of mathematics are either provable in  $\text{RCA}_0$ , or logically equivalent (over  $\text{RCA}_0$ ) to one of  $\text{WKL}_0$ ,  $\text{ACA}_0$ ,  $\text{ATR}_0$ ,  $\Pi_1^1\text{-CA}_0$ .*

**Definition 1.2** The system of **recursive comprehension axioms** ( $\text{RCA}_0$ ) consists of:

- (0) Axioms and inference rules of first-order logic with axioms of equality for numbers.
- (1) Basic arithmetic axioms: Same as  $\text{Q}_<$  (Chapter 4).
- (2)  $\Delta_1^0$  comprehension axiom ( $\Delta_1^0\text{-CA}_0$ ):  $\forall n(\varphi(n) \leftrightarrow \psi(n)) \rightarrow \exists X \forall n(n \in X \leftrightarrow \varphi(n))$ , where  $\varphi(n)$  is  $\Sigma_1^0$ ,  $\psi(n)$  is  $\Pi_1^0$ , and neither includes  $X$  as a free variable.
- (3)  $\Sigma_1^0$  induction:  $\varphi(0) \wedge \forall n(\varphi(n) \rightarrow \varphi(n+1)) \rightarrow \forall n\varphi(n)$ , for any  $\Sigma_1^0$  formula  $\varphi(n)$ .

§2. Real numbers in  $\text{RCA}_0$ 

- $\text{RCA}_0$  is a conservative extension of first-order arithmetic  $\text{I}\Sigma_1$ . (**Lemma 1.3**)
- Primitive recursive functions (e.g., sequence numbers and Gödel numbers) are available in  $\text{RCA}_0$ .

So,  $\mathbb{N}$ ,  $\mathbb{Z}$ ,  $\mathbb{Q}$  and their arithmetical operations are naturally defined in  $\text{RCA}_0$ .

A sequence of rational numbers  $\{q_n\}$  is a **real number**,  $\{q_n\} \in \mathbb{R}$ , if it satisfies

$$\forall n \forall i (|q_n - q_{n+i}| \leq 2^{-n}).$$

The equality  $=$ , the inequality  $<$  and the operations  $+$ ,  $\cdot$  on real numbers are defined as

$$\begin{aligned} \{p_n\} = \{q_n\} &\leftrightarrow \forall n (|p_n - q_n| \leq 2^{-n+1}), & \{p_n\} < \{q_n\} &\leftrightarrow \exists n (q_n - p_n > 2^{-n+1}), \\ \{p_n\} + \{q_n\} &= \{p_{n+1} + q_{n+1}\}, & \{p_n\} \cdot \{q_n\} &= \{p_{n+m} \cdot q_{n+m}\} \text{ for large enough } m. \end{aligned}$$

## Summary

It is provable in  $\text{RCA}_0$  that  $(\mathbb{R}, +, \cdot, 0, 1, <, =)$  is an **Archimedean ordered field**.

# Sequences of real numbers

We define a **sequence of real numbers** as a function  $f : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{Q}$  such that:

for each  $n$ , letting  $f_n(m) = f(n, m)$ ,  $f_n : \mathbb{N} \rightarrow \mathbb{Q}$  is a real number.

The **limit** of  $\{f_n\}$ , denoted  $\lim_{n \rightarrow \infty} f_n$ , is defined as a unique real number  $a$  such that

$$\forall \varepsilon > 0 \exists n \forall i (|a - f_{n+i}| < \varepsilon).$$

## Theorem 2.1 (Nested interval property of $\mathbb{R}$ )

It is provable in  $\text{RCA}_0$  that if  $\{a_n\}$  and  $\{b_n\}$  are two sequences of real numbers such that

$$a_n \leq a_{n+1} \leq b_{n+1} \leq b_n \text{ (for all } n), \text{ and } \lim_{n \rightarrow \infty} |a_n - b_n| = 0,$$

then there exists a real number  $c$  such that  $c = \lim_{n \rightarrow \infty} a_n = \lim_{n \rightarrow \infty} b_n$ .

## Theorem 2.2 (Uncountability of $\mathbb{R}$ )

It is provable in  $\text{RCA}_0$  that for any sequence of real numbers  $\{a_n\}$ , there exists a real number  $c$  such that  $\forall n (a_n \neq c)$ .

We will introduce continuous functions on  $\mathbb{R}$  in  $\text{RCA}_0$ .

### Definition 2.3

A set  $\Phi \subseteq \mathbb{Q}^4$  that satisfies the following conditions is called the **code** for a **continuous function**  $f : \text{dom } f (\subseteq \mathbb{R}) \rightarrow \mathbb{R}$ .

- (1)  $(p, q, r, s) \in \Phi \rightarrow p < q \wedge r \leq s,$
- (2)  $(p, q, r, s), (p', q', r', s') \in \Phi, p' < q \wedge p < q' \rightarrow r' \leq s \wedge r \leq s'.$

Intuitively,  $(p, q, r, s) \in \Phi$  means  $\forall x (p < x < q \rightarrow r \leq f(x) \leq s)$ .

A real number  $x$  belongs to the **domain** of a continuous function  $f$  coded by  $\Phi$ , if

$$\forall n \exists (p, q, r, s) \in \Phi (p < x < q \wedge s - r < 2^{-n}), \text{ denoted } x \in \text{dom } f.$$

It is provable in  $\text{RCA}_0$  that if  $x \in \text{dom } f$ , there exists a unique real  $y$  such that  $\forall (p, q, r, s) \in \Phi (p < x < q \rightarrow r \leq y \leq s)$ . We denote this  $y$  as  $f(x)$ .

### Theorem 2.4 (Intermediate Value Theorem)

The following is provable in  $\text{RCA}_0$ . Given a continuous function  $f$  such that its domain includes  $[0, 1]$  and  $f(0) < 0 < f(1)$ , there exists an  $x \in [0, 1]$  such that  $f(x) = 0$ .

## Definition 3.1

The **system of arithmetical comprehension axioms** ( $ACA_0$ ) is  $RCA_0$  extended with  $(\Pi_0^1\text{-CA}) : \exists X \forall n (n \in X \leftrightarrow \varphi(n))$ , where  $\varphi(n)$  is a  $\Pi_0^1$  (arithmetical) formula with no free occurrence of  $X$ .

## Lemma 3.2

$ACA_0$  is a conservative extension of Peano Arithmetic PA.

## Lemma 3.3

In  $RCA_0$ , the following are equivalent:

- (1)  $ACA_0$ , (2)  $(\Sigma_1^0\text{-CA})$ , (3) The range of any 1-1 function  $f : \mathbb{N} \rightarrow \mathbb{N}$  exists.

## Theorem 3.4

Each of the following propositions is equivalent to  $ACA_0$  in  $RCA_0$ :

- (1) Bolzano-Weierstrass Theorem, (2) Every Cauchy sequence converges,  
(3) Monotone Convergence Theorem: Every bounded increasing sequence converges.

Let  $\text{Seq}_2$  denote the set of (finite) binary sequences. A subset  $T$  of  $\text{Seq}_2$  is called a **tree**, if

$$\forall s, t (s \in \text{Seq}_2 \wedge s \subseteq t \wedge t \in T \implies s \in T),$$

where  $s \subseteq t$  means  $\forall i < \text{length}(s)(s(i) = t(i))$ .

An (infinite) **path** through  $T$  is a function  $g : \mathbb{N} \rightarrow T$  such that  $g(0) = \emptyset$ ,  $g(n) \subseteq g(n+1)$  and  $\neg \exists t (g(n) \subsetneq t \subsetneq g(n+1))$  for all  $n \in \mathbb{N}$ . A path  $g$  is often identified with its range.

### Definition 3.5

**Weak König's lemma** asserts that every infinite subtree of  $\text{Seq}_2$  has an infinite path. The system obtained from  $\text{RCA}_0$  by adding this lemma as an axiom is called  $WKL_0$ .

### Lemma 3.6

In  $\text{RCA}_0$ ,  $WKL_0$  is equivalent to the following statement:

$$(\Sigma_1^0\text{-SP}) : \forall n(\varphi(n) \rightarrow \psi(n)) \rightarrow \exists X \forall n \{(\varphi(n) \rightarrow n \in X) \wedge (n \in X \rightarrow \psi(n))\},$$

where  $\varphi(n)$  is  $\Sigma_1^0$  and  $\psi(n)$  is  $\Pi_1^0$ , neither containing  $X$  as a free variable.

## Corollary 3.7

$WKL_0$  is strictly stronger than  $RCA_0$ .

**Proof** Let  $Rec$  be the set of recursive subsets of  $\omega$ . Then,  $(\omega, Rec)$  is the minimal model of  $RCA_0$ , but it is not a model of  $WKL_0$ , since there are two disjoint  $\Sigma_1^0$  sets  $A$  and  $B$  that are recursively inseparable.  $\square$

However, as shown later,  $WKL_0$  and  $RCA_0$  prove the same arithmetical formulas as  $I\Sigma_1$ , though the first-order part of  $ACA_0$  is just PA. Thus, these three systems have different strength.

Next, we show that the Heine-Borel theorem is equivalent to  $WKL_0$  over  $RCA_0$ . First, we define some basic notions. An **open interval** with rational endpoints  $p, q$  ( $p < q$ ) is represented by the code of the pair  $(p, q)$ . An **open set** of  $\mathbb{R}$  is defined (encoded) as a set of codes of open intervals. We say that an open set  $U$  of  $\mathbb{R}$  **covers** the closed unit interval  $[0, 1]$ , if for any real number  $x \in [0, 1]$ , there exists a code  $(p, q) \in U$  such that  $p < x < q$ .

The **Heine-Borel (Covering) Theorem** asserts that if an open set  $U$  covers the closed interval  $[0, 1]$ , then there exists a finite subset  $U'$  of  $U$  that also covers  $[0, 1]$ .

## Lemma 3.8

The Heine-Borel Theorem is provable in  $WKL_0$ .

**Proof** For each  $s \in \text{Seq}_2$ , we associate an open interval  $(a_s, b_s)$  such that

$$a_s = \sum_{i < \text{leng}(s)} \frac{s(i)}{2^{i+1}}, \quad b_s = a_s + \frac{1}{2^{\text{leng}(s)}}.$$

We can easily see that if  $s \subseteq t$ , then  $[a_t, b_t] \subseteq [a_s, b_s]$ . For any infinite binary decimal  $x = 0.s(0)s(1)s(2)\cdots$ , we have  $x \in [a_{s'}, b_{s'}]$  for all finite segments  $s' \subset s$ .

Now, take any open covering  $U$  of  $[0, 1]$ . For convenience, let's denote the open interval with code  $i$  as  $(p_i, q_i)$ . Then, we define a tree  $T \subseteq \text{Seq}_2$  as follows:

$$\begin{aligned} s \in T &\leftrightarrow \neg \exists i \leq \text{leng}(s) (i \in U \wedge p_i < a_s < b_s < q_i) \\ &\leftrightarrow \forall i \leq \text{leng}(s) (i \in U \rightarrow [a_s, b_s] \not\subseteq (p_i, q_i)). \end{aligned}$$

We first show that  $T$  has no infinite path.

By way of contradiction, we assume that there exists a path  $f \subseteq T$ , where

$$s \in T \leftrightarrow \forall i \leq \text{leng}(s)(i \in U \rightarrow [a_s, b_s] \not\subseteq (p_i, q_i)).$$

By the nested interval property, there exists a (unique) real number  $x$  such that  $a_s \leq x \leq b_s$  for all  $s \in f$ . Since the open set  $U$  covers  $[0, 1]$ , there exists some  $i \in U$  such that the real number  $x$  is contained in the open interval  $(p_i, q_i)$ . Then, there exists an  $s \in f$  with sufficient length such that  $p_i < a_s \leq x \leq b_s < q_i$ , which implies  $s \notin T$ , a contradiction.

Since  $T$  has no infinite path, by weak König's lemma,  $T$  is a finite set. This means that there exists a sufficiently large  $n$  such that all sequences belonging to  $T$  have a length shorter than  $n$ . Thus,

$$\forall s(\text{leng}(s) = n \rightarrow \exists i \leq n(i \in U \wedge p_i < a_s < b_s < q_i)).$$

Then, since  $\cup\{[a_s, b_s] : \text{leng}(s) = n\} = [0, 1]$ ,  $\{i \in U : i \leq n\}$  also covers  $[0, 1]$ . □

## Theorem 3.9

In  $\text{RCA}_0$ , the Heine-Borel Theorem is equivalent to  $\text{WKL}_0$ .

**Proof** We have already shown that the Heine-Borel Theorem holds in  $\text{WKL}_0$ . Now, we assume the Heine-Borel Theorem and derive the weak König's lemma.

First, let's discuss the idea behind the proof. The Heine-Borel Theorem implies the compactness of  $[0, 1]$ , which leads to the compactness of the ternary set (a closed set)

$$\left\{ \sum_{i=0}^{\infty} \frac{f(i)}{3^{i+1}} \mid f \in \{0, 2\}^{\mathbb{N}} \right\},$$

and hence also the compactness of the Cantor space  $\{0, 1\}^{\mathbb{N}}$  since it is homeomorphic to the ternary set. Finally, the compactness of  $\{0, 1\}^{\mathbb{N}}$  implies  $\text{WKL}_0$ .

For preparation, for each  $s \in \text{Seq}_2$ , we associate the rational interval  $(a_s, b_s)$  defined as:

$$a_s = \sum_{i < \text{leng}(s)} \frac{2s(i)}{3^{i+1}}, \quad b_s = a_s + \frac{1}{3^{\text{leng}(s)}}.$$

Then, the closed intervals  $[a_{s \cap 0}, b_{s \cap 0}]$  and  $[a_{s \cap 1}, b_{s \cap 1}]$  respectively become the left and right thirds of the interval  $[a_s, b_s]$ . Here,  $s \cap i$  denotes the sequence  $s$  followed by  $i = 0, 1$ . So, for any real number  $x$  not belonging to the ternary set, there exists exactly one open interval  $(b_{s \cap 0}, a_{s \cap 1})$  containing it. Hence,

$$\bigcup \{(b_{s \cap 0}, a_{s \cap 1}) \mid s \in \text{Seq}_2\}$$

is the complement of the ternary set.

Furthermore, for each  $s \in \text{Seq}_2$ , define

$$a'_s = a_s - \frac{1}{3^{\text{leng}(s)+1}}, \quad b'_s = b_s + \frac{1}{3^{\text{leng}(s)+1}}.$$

Then, we have  $[a_s, b_s] \subset (a'_s, b'_s)$  for each  $s \in \text{Seq}_2$ , but  $(a'_s, b'_s)$  and  $(a'_t, b'_t)$  intersect only if either  $s$  or  $t$  is an initial segment of the other. For any real number  $x$  in the ternary set, there exists a unique  $f \in \{0, 1\}^{\mathbb{N}}$  such that  $x \in (a'_s, b'_s)$  for any finite initial sequence  $s \subset f$ .

Now, to prove the weak König's lemma, take any (nonempty) tree  $T \subseteq \text{Seq}_2$  without infinite paths and show that  $T$  is finite.

Let  $B$  be the set of minimal binary sequences not in  $T$ , that is,

$$s \in B \Leftrightarrow s \notin T \wedge \forall t \subset s (t \neq s \rightarrow t \in T).$$

It's clear that any infinite path  $f \subseteq \text{Seq}_2$  shares exactly one element  $s \in B$  and  $s \subset f$ .

Then, the following set  $U$  covers  $[0, 1]$ :

$$U = \bigcup \{(a'_s, b'_s) : s \in B\} \cup \bigcup \{(b_{s \cap 0}, a_{s \cap 1}) : s \in \text{Seq}_2\}.$$

So, by the Heine-Borel Theorem,  $U$  has a finite subcover  $U'$ .

Since for any  $s \in B$ ,  $(a'_s, b'_s)$  does not intersect with any other  $(a'_t, b'_t) \in U$  and is not a subset of  $\bigcup \{(b_{s \cap 0}, a_{s \cap 1}) : s \in \text{Seq}_2\}$ ,  $U'$  must contain  $\{(a'_s, b'_s) : s \in B\}$ . Therefore,  $B$  is finite.

Since  $T$  is obtained from the set of all initial segments of elements in  $B$  by removing the elements of  $B$ , it is also finite. □

The Heine-Borel property of  $[0, 1]$  allows us to derive various properties of continuous functions  $f : [0, 1] \rightarrow \mathbb{R}$ .

### Lemma 3.10

In  $\text{WKL}_0$ , a continuous function  $f : [0, 1] \rightarrow \mathbb{R}$  is uniformly continuous.

**Proof** Fix any  $n \in \mathbb{N}$ . We want to show the existence of  $d > 0$  such that

$$\forall x, y \in [0, 1] (|x - y| < d \rightarrow |f(x) - f(y)| < 2^{-n}).$$

Let  $F$  be the code for the continuous function  $f$ , and denote the open interval with code  $i$  as  $(p_i, q_i)$ . Then, define the open set  $U$  as follows:

$$i \in U \Leftrightarrow \exists j < i ((p_i, q_i, p_j, q_j) \in F \wedge q_j - p_j < 2^{-n-1}).$$

First, we show that  $U$  is a covering of  $[0, 1]$ . For any real number  $x \in [0, 1]$ , since  $x \in \text{dom} f$ , there exists  $(p_k, q_k, p_j, q_j) \in F$ , such that

$$p_k < x < q_k \wedge q_j - p_j < 2^{-n-1}.$$

Furthermore, there are infinitely many  $i$  such that  $p_k \leq p_i < x < q_i \leq q_k$ , so taking such an  $i > j$ , we have  $i \in U$  with  $p_i < x < q_i$ . Therefore,  $U$  forms an open covering of  $[0, 1]$ .

By the Heine-Borel Theorem,  $U$  has a finite subcover  $U'$ .

Let  $d$  be the minimum width  $q_i - p_i$  among the intervals  $(p_i, q_i)$  in  $U'$ . We shall show that this  $d$  satisfies the uniform convergence condition.

Now, choose any real numbers  $x, y \in [0, 1]$  such that  $0 < y - x < d$ . Then, there must exist intervals  $(p_i, q_i), (p_{i'}, q_{i'})$  in  $U'$  such that  $x \in (p_i, q_i)$ ,  $y \in (p_{i'}, q_{i'})$  and that they have a common point  $z$ .

By way of contradiction, we assume that there are no such intervals. Then, we take an interval  $(p_i, q_i) \ni x$  in  $U'$  with maximal  $q_i$ , and an interval  $(p_{i'}, q_{i'}) \ni y$  in  $U'$  with minimal  $p_{i'}$ , and still have  $q_i < p_{i'}$ . Since  $U'$  is a covering, there exists  $(p_k, q_k) \ni q_i$  in  $U'$ . By the maximality of  $q_i$ ,  $x \notin (p_k, q_k)$ . Since  $|q_k - p_k| \geq d > |x - y|$ , we have  $y \in (p_k, q_k)$ , which contradicts with the minimality of  $p_{i'}$ .

By the definition of  $U$ , we have  $|f(x) - f(z)| < 2^{-n-1}$  and  $|f(y) - f(z)| < 2^{-n-1}$ , thus  $|f(x) - f(y)| < 2^{-n}$ , which fulfills the lemma.  $\square$

## Lemma 3.11

In  $\text{WKL}_0$ , a continuous function  $f : [0, 1] \rightarrow \mathbb{R}$  attains a maximum value.

**Proof** First, we show that the supremum  $M$  of the range of  $f$  exists.

As in the proof of the previous lemma, we define  $U$  by a  $\Sigma_0^0$  formula:

$$i \in U \Leftrightarrow \exists j < i ((p_i, q_i, p_j, q_j) \in F \wedge q_j - p_j < 2^{-n-1}).$$

We can finitely calculate whether or not a given finite set of open rational intervals covers  $[0, 1]$ . Therefore, in  $\text{WKL}_0$ , we can construct a function extracting  $U'$  according to  $n$ .

For each  $i \in U'$ , select  $j_i < i$  such that  $(p_i, q_i, p_{j_i}, q_{j_i}) \in F \wedge q_{j_i} - p_{j_i} < 2^{-n-1}$ , and let  $M_n = \max\{q_{j_i} : i \in U'\}$ . Thus, we can construct a sequence  $\{M_n\}$ , which becomes a real number, and it is clear that it is the supremum  $M$  of the range of  $f$ .

What remains is to show that the existence of a point  $x = a$  such that  $f(a) = M$ . For the sake of the following argument, we redefine  $M_n = \max\{p_{j_i} : i \in U'\}$ . This ensures that for any  $n$ ,  $M_n \leq M = \{M_n\}$ .

By way of contradiction, assume that  $f(x) < M$  for all  $x \in [0, 1]$ . Then, we define an open set  $V$  as follows:

$$i \in V \Leftrightarrow \exists j < i \exists n < i ((p_i, q_i, p_j, q_j) \in F \wedge q_j < M_n).$$

To show that this set forms a covering of  $[0, 1]$ , take any real number  $x \in [0, 1]$ . Since  $f(x) < M$ , there exists  $n$  such that  $f(x) < M_n \leq M$ , and hence there exists  $(p_k, q_k, p_j, q_j) \in F$  and  $n$  such that

$$p_k < x < q_k \wedge p_j \leq f(x) \leq q_j < M_n \leq M.$$

As there are infinitely many  $i$  such that  $p_k \leq p_i < x < q_i \leq q_k$ , taking  $i > \max\{j, n\}$  ensures  $i \in V$  with  $p_i < x < q_i$ . Therefore,  $V$  forms an open covering of  $[0, 1]$ .

Again, by the Heine-Borel Theorem,  $V$  has a finite subcover  $V'$ . Let  $M'$  be the maximum of  $q_i$  for  $(p_i, q_i)$  in  $V'$ . Then, by the definition of values of a continuous function, obviously  $M'$  is an upper bound of the range. However, due to the finiteness of  $V'$ , for some  $n$ ,  $M' < M_n \leq M$ , which contradicts the fact that  $M$  is the supremum.  $\square$

Conversely, the properties described in the two lemmas above allow us to derive  $WKL_0$ . In sum, the following theorem holds:

### Theorem 3.12

The following assertions are pairwise equivalent in  $RCA_0$ :

- (1)  $WKL_0$ ,
- (2) A continuous function  $f : [0, 1] \rightarrow \mathbb{R}$  is uniformly continuous,
- (3) A continuous function  $f : [0, 1] \rightarrow \mathbb{R}$  is bounded,
- (4) A bounded continuous function  $f : [0, 1] \rightarrow \mathbb{R}$  has a supremum,
- (5) A continuous function  $f : [0, 1] \rightarrow \mathbb{R}$  that has a supremum attains its maximum value.

**Proof** By Lemmas 3.10 and 3.11, we can derive (2), (3), (4), and (5) from (1). Hence, it suffices to obtain counterexamples for (2), (3), (4) and (5) from the negation of (1). Now, assume the negation of (1). Then, there exists an infinite tree  $T \subseteq \text{Seq}_2$  without infinite paths.

As shown in the proof of Heine-Borel's theorem, for each  $s \in \text{Seq}_2$ , define the two rational numbers  $a_s$  and  $b_s$  as follows:

$$a_s = \sum_{i < \text{leng}(s)} \frac{s(i)}{2^{i+1}},$$

$$b_s = a_s + \frac{1}{2^{\text{leng}(s)}}.$$

Let  $B$  be the infinite set of all minimal binary sequences not in  $T$ ,

$$s \in B \Leftrightarrow s \notin T \wedge \forall t \subset s (t \neq s \rightarrow t \in T)$$

and  $J$  be the set of closed intervals  $[a_s, b_s]$  for all  $s \in B$ .

Each real number  $x \in [0, 1]$  is either an interior point of exactly one interval in  $J$  or an endpoint of one or two intervals. Such an infinite set  $J$  is called a **singular closed cover**.

$\neg \text{WKL}_0 \rightarrow \neg (3)$  bounded,  $\neg (5)$  a maximum value.

We will construct a counterexample for (3) using this singular closed cover  $J$ . This also serves as a counterexample for (2) since (2) implies (3). We define a continuous function  $f_s$  for each interval  $[a_s, b_s]$  in  $J$  as follows:

$$f_s(x) = \begin{cases} \text{leng}(s) \frac{2(x-a_s)}{a_s+b_s} & \text{if } a_s \leq x \leq \frac{a_s+b_s}{2}, \\ \text{leng}(s) \frac{2(b_s-x)}{a_s+b_s} & \text{if } \frac{a_s+b_s}{2} \leq x \leq b_s. \end{cases}$$

That is,  $f_s$  takes 0 at the endpoints  $x = a_s, b_s$ , takes  $\text{leng}(s)$  at the midpoint  $x = \frac{a_s+b_s}{2}$ , and is linearly interpolated otherwise.

Let  $f$  be a function obtained by composing all such functions  $f_s$ . Then, it is clearly continuous but unbounded. (Exercise: construct a continuous function code for  $f$ .)

A counterexample for (5) can be constructed in the similar way as follows:

$$f_s(x) = \begin{cases} (1 - 2^{-\text{leng}(s)}) \frac{2(x-a_s)}{a_s+b_s} & \text{if } a_s \leq x \leq \frac{a_s+b_s}{2}, \\ (1 - 2^{-\text{leng}(s)}) \frac{2(b_s-x)}{a_s+b_s} & \text{if } \frac{a_s+b_s}{2} \leq x \leq b_s. \end{cases}$$

Then,  $f$  clearly has 1 as its supremum, but it can not attain the maximum value 1 in  $[0, 1]$ .

$\neg \text{WKL}_0 \rightarrow \neg (4)$  a supremum.

Theorem 3.4.

$(\text{RCA}_0 \vdash) \text{ACA}_0 \Leftrightarrow (5)$  Every bounded increasing sequence of reals has a supremum.

Negating  $\text{WKL}_0$ , we have the negation of  $\text{ACA}_0$ , which implies the existence of a bounded increasing sequence of rational numbers  $\{c_n\}$  that lacks a supremum.

Then, replace the maximum value of  $f_s$  with  $c_{\text{leng}(s)}$  and proceed similarly.

$$f_s(x) = \begin{cases} c_{\text{leng}(s)} \frac{2(x-a_s)}{a_s+b_s} & \text{if } a_s \leq x \leq \frac{a_s+b_s}{2}, \\ c_{\text{leng}(s)} \frac{2(b_s-x)}{a_s+b_s} & \text{if } \frac{a_s+b_s}{2} \leq x \leq b_s, \end{cases}$$

HW Problem # 5-3

Show that in the theorem 3.12 (4) and (5), "continuous function" can be replaced with "uniformly continuous function". Hint: It is beneficial to use a singular closed cover for the ternary set.

- The deadline for the fifth homework is changed to December 10th.
- The final report problems will be assigned on Dec. 10 and 12, and due Jan. 12.

# Complete separable metric spaces and continuous functions

Let  $A$  be a non-empty subset of  $\mathbb{N}$ . Suppose  $d : A \times A \rightarrow \mathbb{R}$  is a (pseudo) metric on  $A$ . A sequence  $\{a_n\}$  from  $A$  satisfying  $\forall n \forall i d(a_n, a_{n+i}) \leq 2^{-n}$  is called a point of  $\hat{A}$ , and we write  $\{a_n\} \in \hat{A}$ .  $\hat{A}$  can be viewed as a **complete separable metric space**.

**Example.** If  $A = \mathbb{Q}$  and  $d(p, q) = |p - q|$ , then  $\hat{A}$  is nothing but  $\mathbb{R}$ . Also, if  $A = \mathbb{Q}^2$  and  $d((p, q), (p', q')) = \sqrt{(p - p')^2 + (q - q')^2}$ , then  $\hat{A}$  is  $\mathbb{R}^2$ .

In a metric space  $\hat{A}$ , an **open ball**  $B_r(a)$  centered at  $a \in A$  with a rational radius  $r > 0$  is coded by the pair  $(a, r) (\in A \times \mathbb{Q}^+)$ . An **open set** is a set of codes of open balls.

The code  $F$  of a **continuous function**  $f$  from a metric space  $\hat{A}$  to a metric space  $\hat{B}$  is a subset of  $A \times \mathbb{Q}^+ \times B \times \mathbb{Q}^+$ , fulfilling conditions similar to those for a continuous function from  $\mathbb{R}$  to  $\mathbb{R}$ , by which

$$(a, r, b, s) \in F \text{ means } x \in B_r(a) \rightarrow f(x) \in \overline{B_s(b)} \text{ (closed ball).}$$

## Brouwer's Fixed-Point Theorem

**Brouwer's Fixed-Point Theorem** states that any continuous function  $f : [0, 1]^n \rightarrow [0, 1]^n$  has a fixed point, i.e., a point  $x$  such that  $f(x) = x$ .

## Theorem 3.13 (Shioji-T.)

Brouwer's Fixed-Point Theorem is equivalent to  $WKL_0$  over  $RCA_0$ .

In  $WKL_0$ , Brouwer's Fixed-Point Theorem can be extended to the infinite-dimensional space  $[0, 1]^{\mathbb{N}} (\subseteq \mathbb{R}^{\mathbb{N}})$ , which is known as the **Tychonoff-Schauder fixed-point theorem**.

By utilizing this fixed-point theorem, the **Cauchy-Peano theorem** for the existence of local solutions to ordinary differential equations can be derived within  $WKL_0$ , and the converse is also provable.

Various fixed-point theorems and their applications (e.g., the Hahn-Banach theorem) have been studied by N. Shioji and K. Tanaka [Fixed point theory in weak second-order arithmetic *Ann. Pure Appl. Logic*, 47, 167-188, 1990].

## §4. König's Lemma and Ramsey's theorem

We begin with König's Lemma, not the "weak" version.

Let  $\text{Seq}$  denote the set of finite sequences from  $\mathbb{N}$ , that is, the set of functions with domain  $\{i \in \mathbb{N} : i < n\}$  for some  $n \in \mathbb{N}$ .

A subset  $T$  of  $\text{Seq}$ , which is closed under initial segment, is called a **tree**.

A tree  $T$  is said to be **finitely branching**, if each node  $s \in T$  has at most finitely many children, i.e.,

$$\forall s(s \in T \rightarrow \exists n \forall m(s \frown m \in T \rightarrow m < n))$$

A subtree of  $T$  that never branches is called a **path** of  $T$ .

**König's Lemma** asserts that "every infinite, finitely branching tree has an infinite path." Weak König's Lemma is König's Lemma about special trees consisting of binary sequences. As we will see, König's Lemma is equivalent to  $\text{ACA}_0$ , and thus it is properly stronger than weak König's Lemma.

## Theorem 4.1

Over  $\text{RCA}_0$ , the following are pairwise equivalent:

- (1)  $\text{ACA}_0$
- (2) König's Lemma
- (3) An infinite tree  $T$ , such that each node  $s \in T$  has at most two children  $s \frown m \in T$  ( $m \in \mathbb{N}$ ), has an infinite path.

Note: In the above (3), it is crucial that  $m$  such that  $s \frown m \in T$  is not bounded over  $T$ . If  $m$  were bounded, the assertion would be equivalent to weak König's Lemma.

**Proof** (1)  $\Rightarrow$  (2). Given an infinite finitely branching tree  $T$ , let  $T'$  be the set of  $s \in T$  that have an infinitely many descendants  $t \supseteq s$  (by  $(\Pi_0^1\text{-CA})$ ).

Then, using primitive recursion, define a path  $g$  in  $T'$  as follows:

$$g(0) = \text{empty sequence,}$$

$$g(n+1) = g(n) \frown m, \text{ where } m \text{ is the smallest number such that } g(n) \frown m \in T'.$$

(2)  $\Rightarrow$  (3) is trivial. To show (3)  $\Rightarrow$  (1), assume (3) and show the existence of range of a given 1-1 function  $f : \mathbb{N} \rightarrow \mathbb{N}$ , which is equivalent to  $\text{ACA}_0$ , by Lemma 3.3.(3).

Define a tree  $T$  as follows:  $s \in T \Leftrightarrow$

$$(a) \quad \forall m, n < \text{leng}(s) (f(m) = n \leftrightarrow s(n) = m + 1),$$

$$(b) \quad \forall n < \text{leng}(s) (s(n) > 0 \rightarrow f(s(n) - 1) = n).$$

Then, each node  $t \in T$  has at most two children  $t \frown k \in T$ . This is because letting  $s = t \frown k$ ,  $n = \text{leng}(s) - 1$  in (b), we have  $k = 0$  or  $k = f^{-1}(n) + 1$  if  $f^{-1}(n)$  exists.

Next, show that  $T$  is infinite. For this, it suffices to show that for any  $k \in \mathbb{N}$ , there exists a sequence  $s \in T$  with  $\text{leng}(s) = k$ . First, by bounded  $(\Sigma_1^0\text{-CA})$ ,  $Y = \text{ran} f \cap k$ , that is,  $\{n \in \text{ran} f : n < k\}$  exists. Then, define a sequence  $s$  of length  $k$  as follows: for  $n < k$ ,

$$s(n) = \begin{cases} 0 & \text{if } n \notin Y \\ m + 1 & \text{if } n \in Y \wedge f(m) = n \end{cases}$$

Obviously,  $s \in T$ . So,  $T$  satisfies the conditions of (3).

Now, by (3), the tree  $T$  has an infinite path  $g$ . From the condition (a) of  $T$ ,

$$\forall m, n (f(m) = n \leftrightarrow g(n) = m + 1).$$

Thus, setting  $X = \{n : g(n) > 0\}$ , we have  $X = \text{ran} f$ . □

**Ramsey's Theorem** was first invented by F. Ramsey in order to settle Hilbert's decision problem for first-order logic, though he only succeeded partially and we have no space to explain his original motivation and results.

For a set  $X \subseteq \mathbb{N}$ , we denote by  $[X]^k$  the set of all sequences  $(m_1, \dots, m_k)$  of  $k$  elements from  $X$  such that  $m_1 < \dots < m_k$ . Somewhat naïvely, (infinite) Ramsey's theorem  $\text{RT}_l^k$  states that for a coloring of  $[\mathbb{N}]^k$  into  $l$  colors, there exists an infinite subset  $X \subseteq \mathbb{N}$  such that  $[X]^k$  is monochromatic<sup>1</sup>. More precisely, we state it as follows.

### Definition 4.2 (Ramsey's Theorem)

Let  $k, l > 0$  be natural numbers. Ramsey's Theorem  $\text{RT}_l^k$  is the following assertion:

$$\forall f : [\mathbb{N}]^k \rightarrow \{0, 1, \dots, l - 1\} \exists X \subseteq \mathbb{N} (X \text{ is infinite} \wedge f \text{ is constant on } [X]^k).$$

---

<sup>1</sup>Finite Ramsey's Theorem, denoted  $m \rightarrow (n)_l^k$ , is the statement that if  $\{0, \dots, n - 1\}^k$  is painted in  $l$  colors, there exists a subset  $X \subseteq \{0, \dots, n - 1\}$  of  $m$  elements such that  $[X]^k$  is monochromatic. The finite version can be derived from the infinite version by the compactness argument.

For example,  $\text{RT}_l^2$  can be interpreted as follows: If all pairs  $\{m, n\}$  of natural numbers are painted in  $l$  colors, then there always exists an infinite set  $X$  such that all pairs of elements from  $X$  are painted the same color. Such an  $X$  is called a **homogeneous set**.

If we consider the statement of painting any finite number of colors, we denote it as  $\text{RT}^k$ , i.e.,  $\text{RT}^k \equiv \forall l \in \mathbb{N}(\text{RT}_l^k)$ .

Although  $\text{RT}_l^k$  for any standard natural number  $l \geq 2$  can be deduced from  $\text{RT}_2^k$  by meta-induction in  $\text{RCA}_0$ , the equivalence of  $\text{RT}^k$  to  $\text{RT}_2^k$  may require  $\Pi_2^1$ -induction, since  $\text{RT}_l^k$  is a  $\Pi_2^1$  formula.

First consider the strength of  $\text{RT}^1$ , which is also known as the (infinite version of) **pigeonhole principle** (PHP). For a standard natural number  $l \geq 1$ ,  $\text{RT}_l^1$  obviously holds even in  $\text{RCA}_0$ . The question is how much restricted induction is needed to derive  $\forall l \text{RT}_l^1$ .

Recall: the collection principle  $(\text{B}\varphi)$  for  $\varphi(x, y_1, \dots, y_k)$  in  $\mathcal{L}_{OR}$  is as follows

$$\forall x < u \exists y_1 \cdots \exists y_k \varphi(x, y_1, \dots, y_k) \rightarrow \exists v \forall x < u \exists y_1 < v \cdots \exists y_k < v \varphi(x, y_1, \dots, y_k).$$

$\text{B}\Pi_1^0$  denotes  $\{(\text{B}\varphi) \mid \varphi \in \Pi_1^0\}$ .  $\text{B}\Pi_1^0$  is equivalent to  $\text{B}\Sigma_2^0$ , and  $\text{I}\Sigma_1^0 \subsetneq \text{B}\Sigma_2^0 \subsetneq \text{I}\Sigma_2^0$ .  $\text{B}\Pi_1^0$  is not provable in  $\text{WKL}_0$ , but obviously provable in  $\text{ACA}_0$ .

## Theorem 4.3 (J. Hirst)

In  $\text{RCA}_0$ ,  $\text{RT}^1$  is equivalent to  $\text{BII}_1^0$ .

The above theorem indicates that the strength of  $\text{RT}^1$  is intermediate between  $\text{ACA}_0$  and  $\text{RCA}_0$ , and it is incomparable with  $\text{WKL}_0$ .

## Theorem 4.4

In  $\text{ACA}_0$ , both  $\text{RT}^1$  and  $\forall k(\text{RT}^k \rightarrow \text{RT}^{k+1})$  are provable.

Since  $\text{RT}^k$  is a  $\text{II}_2^1$  statement, the theorem above does not allow us to derive  $\forall k \text{RT}^k$  within induction of  $\text{ACA}_0$ . Paris and Harrington formulated a proposition PH in first-order arithmetic to express something like  $\forall k \text{RT}^k$ , and proved that PH is independent from PA.

## Lemma 4.5

Within  $\text{RCA}_0$ ,  $\text{ACA}_0$  can be derived from  $\text{RT}_2^3$ .

## Theorem 4.6

For any standard numbers  $k \geq 3$ ,  $l \geq 2$ ,  $\text{RT}_l^k$ ,  $\text{RT}^k$ , and  $\text{ACA}_0$  are equivalent over  $\text{RCA}_0$ .

Finally, concerning  $\text{RT}^2$  and  $\text{RT}_2^2$ , it is known that both are between  $\text{ACA}_0$  and  $\text{RCA}_0$ , and are incomparable with  $\text{WKL}_0$ . Within  $\text{RCA}_0$ ,  $\text{RT}^2$  implies  $\text{BII}_2^0$ , but  $\text{RT}_2^2$  does not.

Thank you for your attention!