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### Logic and Foundation I Part 2. First-order logic

Kazuyuki Tanaka

BIMSA

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### Logic and Foundations I -

- Part 1. Equational theory
- Part 2. First order theory
- Part 3. Model theory
- Part 4. First order arithmetic and incompleteness theorems

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Part 2. Schedule

- Nov. 16, (1) ∀-theory and ∀∃-theory
- Nov. 23, (2) Horn theory and reduced products
- Nov. 30, (3) Ultra products and non-standard analysis

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Today's topics

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### <span id="page-3-0"></span>Definition

Let I be a non-empty set.  $\mathcal{F} \subseteq \mathcal{P}(I)$  is said to be filter on I if the following are satisfied.  $(1) \varnothing \notin \mathcal{F}, I \in \mathcal{F}.$ (2)  $X \in \mathcal{F}, X \subseteq Y \subseteq I \Rightarrow Y \in \mathcal{F}.$ 

(3)  $X, Y \in \mathcal{F} \to X \cap Y \in \mathcal{F}$ .

### Let *I* be an infinite set.

**1** The collection of co-finite subsets of  $I$  is a filter, called a **Fréchet filter**.

**2** For each  $i \in I$ ,  $\{X \subseteq I : i \in X\}$  is a filter, called a **principal filter**.

### Lemma

If  $S \subset \mathcal{P}(I)$  has the finite intersection property: for any finite subset  $\{J_1, \ldots, J_n\} \subset S$ ,

$$
J_1 \cap \cdots \cap J_n \neq \varnothing,
$$

then there exists a filter  $F$  including  $S$ .

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### Definition

Let  $\mathfrak{A}_i = (A_i, \mathtt{f}^{\mathfrak{A}_i}, \ldots, \mathtt{R}^{\mathfrak{A}_i}, \ldots)$   $(i \in I)$  be an  $\mathcal L$ -structure. Let  ${\mathcal F}$  be a filter on  $I.$  Then, we define the binary relation  $\approx_{\mathcal F}$  on  $\prod A_i$  as follows

$$
a \approx_{\mathcal{F}} b \quad \Leftrightarrow \quad \{i \in I : a(i) = b(i)\} \in \mathcal{F}.
$$

### Lemma

### $\approx$   $\tau$  is a congruence relation.

Thus, we can define the quotient structure in the same way as for the algebraic structure.

### Definition

Let  $\mathfrak{A}_i = (A_i, \mathtt{f}^{\mathfrak{A}_i}, \ldots, \mathtt{R}^{\mathfrak{A}_i}, \ldots)$   $(i \in I)$  be  $\mathcal L$ -structures. Let  $\mathcal F$  be a filter on  $I$ . Then, the following  $\mathcal L$ -structure is called the reduced product of  $\mathfrak{A}_i$ , denoted by  $\prod \mathfrak{A}_i/\mathcal F.$ 

$$
\left(\prod A_i/\mathcal{F},\mathbf{f}^{\prod \mathfrak{A}_i/\mathcal{F}},\ldots,\mathrm{R}^{\prod \mathfrak{A}_i/\mathcal{F}},\ldots\right)
$$

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- $\theta_0 \vee \neg \theta_1 \vee \cdots \vee \neg \theta_n$  and  $\neg \theta_1 \vee \cdots \vee \neg \theta_n$  are called **basic Horn formulas**, if  $\theta_i$  ( $i < n$ ) are atomic formulas.
- A formula constructed from the basic Horn formulas by using only ∧, ∀, and ∃ is called a **Horn formula**. A set of Horn sentences is called a **Horn theory**.

For  $a_1, \ldots, a_n \in \prod A_i$ , we set  $\|\varphi(a_1, \ldots, a_n)\| := \{i \in I : \mathfrak{A}_i \models \varphi(a_1(i), \ldots, a_n(i))\}.$ 

### Lemma

Let  $\varphi(x_1,\ldots,x_n)$  be a Horn formula, then for  $a_1,\ldots,a_n\in \prod A_i$ ,

$$
\|\varphi(a_1,\ldots,a_n)\| \in \mathcal{F} \Rightarrow \prod \mathfrak{A}_i/\mathcal{F} \models \varphi([a_1],\ldots,[a_n]).
$$

### Theorem (Keisler-Galvin)

The following are equivalent:

- $(1)$   $Mod(T)$  is closed under reduced products.
- (2) There exists a Horn theory T' such that  $Mod(T) = Mod(T')$ .

A proof (1)  $\Rightarrow$  (2) can be found in Chang-Keisler's classic textbook Model Theory.

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- A sentence with several  $\forall$  in front of a basic Horn formula is called a  $\forall$ -**Horn** sentence (or simply called a Horn sentence in some literature). A collection of such sentences is called a  $\forall$ -**Horn theory** (or simply a Horn theory).
- A ∀-Horn theory is a nice extension of equational theory. The following theorem is a counter part of Birkhoff's equational class theorem. It can be proven similarly, and we leave the details to the reader.

### Theorem

Let  $K$  be a class of  $L$ -structures, then the following are equivalent:

- $(1)$  K is closed under direct products, substructures, and isomorphic images.
- $(2)$  K is closed under reduced products, substructures, and isomorphic images.
- (3) There exists a  $\forall$ -Horn theory T such that  $Mod(T) = K$ .

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### <span id="page-7-0"></span>**Ultraproducts** In the following, we will consider the necessary and sufficient conditions for a class of structures to be axiomatized by first order logic, that is, be expressed as  $Mod(T)$ .

### Definition

A class K of  $\mathcal L$ -structures is called an **elementary class** if there exists a set T of sentences such that  $K = Mod(T)$ . In this case, we write

 $\mathcal{K} \in EC_{\Delta}$ .

To characterize elementary classes, we use a kind of reduced product called "ultraproduct". To define it, we first introduce an ultrafilter.

### Definition

The filter F on I is an **ultrafilter** (maximal filter) if the following properties are satisfied.

 $\forall X \subset I(X \in \mathcal{F} \lor I - X \in \mathcal{F}).$ 

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### Lemma

Every filter  $F$  can be expanded to an ultrafilter  $U$ .

**Proof.** Consider the class of all filters including a given filter  $\mathcal{F}$ . Since it is closed under the union of chains, by Zorn's lemma, there is a maximal filter  $U$ which is an ultrafilter.

A principal filter is an ultrafilter. There exists an ultrafilter which is non-principal.

### Lemma

There exists a non-principal ultrafilter  $U$  on any infinite set  $I$ .

### Proof.

Let I be an infinite set, and F be a Fréchet filter on it (a subset of I whose complement is finite). By the above lemma, an ultrafilter U can be obtained by expanding  $\mathcal F$ . Then U is non-principal, since for each  $i \in I$ ,  $I - \{i\} \in \mathcal{F} \subseteq \mathcal{U}$ , so we have  $\{i\} \notin \mathcal{U}$ .

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## Stone's representation theorem

We prove Stone's representation theorem using an ultrafilter.

### Theorem (Stone's representation theorem)

For any Boolean algebra  $\mathfrak{B}$ , there exists a set X, and  $\mathfrak{B}$  can be embedded in the power set algebra  $\mathfrak{P}(X)$ . In particular, if  $\mathfrak{B}$  is finite, it is isomorphic to  $\mathfrak{P}(X)$ .

### Proof.

- Let  $\mathfrak{B} = (B, \vee, \wedge, \neg, 0, 1)$  be a Boolean algebra. Filters, Ultrafilters, and others can naturally be defined for a subset  $F \subseteq B$  with the ordering  $x \le y \Leftrightarrow x \wedge y = x$ . Let X be the set of all ultrafilters of B and  $\mathcal{P}(X)$  be its power set.
- Define  $f : B \to \mathcal{P}(X)$  as follows:  $f(b)$  is the set of ultrafilters containing b. Then,  $f : B \to \mathcal{P}(X)$  is embedding.
- If  $\mathfrak B$  is finite, any ultrafilter must be a principal filter. And its generator is an atom (non-zero minimal element) in  $\mathfrak{B}$ . So, let X be the set of atoms. It is easy to see that  $\mathfrak{B}$  and  $\mathfrak{B}(X)$  are isomorphic.

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## **Ultraproducts**

### Definition (Ultraproduct)

The reduced product  $\prod \mathfrak{A}_i/\mathcal{U}$  for an ultrafilter  $\mathcal U$  is called an ultraproduct.

### Theorem (Łos)

Let U be an ultrafilter. For any formula  $\varphi(x_1,\ldots,x_n)$  and  $a_1,\ldots,a_n \in \prod A_i$ ,<br> $\prod \mathfrak{A}_i/\mathcal{U} \models \varphi([a_1],\ldots,[a_n]) \Leftrightarrow \|\varphi(a_1,\ldots,a_n)\| \in \mathcal{U}$ .  $\prod \mathfrak{A}_i/\mathcal{U} \models \varphi([a_1], \ldots, [a_n]) \Leftrightarrow ||\varphi(a_1, \ldots, a_n)|| \in \mathcal{U}.$ 

Proof. By induction on the construction of formulas. The atomic formulas and formulas beginning with  $\wedge$  and  $\exists$  can be treated in the same way as reduced products. Then we only need to treat the case of negation  $\neg \varphi$ .

$$
\prod \mathfrak{A}_i/\mathcal{U} \models \neg \varphi \Leftrightarrow \prod \mathfrak{A}_i/\mathcal{U} \not\models \varphi
$$
  
\n
$$
\Leftrightarrow ||\varphi|| \not\in \mathcal{U} \quad (\because \text{ induction hypothesis})
$$
  
\n
$$
\Leftrightarrow ||\neg \varphi|| \in \mathcal{U} \quad (\because \text{maximality of } \mathcal{U}).
$$

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### <span id="page-11-0"></span> $\sim$ Problem 9  $\sim$

Use ultraproducts to show that any field  $\mathcal F$  has algebraic closure  $\overline{\mathcal F}$ .

- We fix a field  $F$  in a language with constants for their elements.
- Let  $\mathcal{F}_P$  be a splitting field of a polynomial P, and for each  $Q \in \mathcal{F}[X]$ , we put

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• Then,  $\{J_{Q}: Q \in \mathcal{F}[X] \}$  and  $Q$  is not a constant. I has the finite intersection property

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### $\sim$ Problem 9  $\sim$

Use ultraproducts to show that any field  $\mathcal F$  has algebraic closure  $\overline{\mathcal F}$ .

### Solution:

- We fix a field  $\mathcal F$  in a language with constants for their elements.
- Let  $\mathcal{F}_P$  be a splitting field of a polynomial P, and for each  $Q \in \mathcal{F}[X]$ , we put

 $J_Q = \{P \in \mathcal{F}[X] : Q$  is splitted into linear factors over  $\mathcal{F}_P\}.$ 

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• Then,  $\{J_Q: Q \in \mathcal{F}[X]$  and  $Q$  is not a constant.} has the finite intersection property  $(\because Q_1 \cdots Q_n \in J_{Q_1} \cap \cdots \cap J_{Q_n})$  . Therefore, it can be expanded to an ultrafilter  $\mathcal U$ .

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- Now consider the ultraproduct  $\prod \mathcal{F}_P/\mathcal{U}$ , which is a field extension of  $\mathcal{F}.$
- For any (non-constant) polynomial  $Q \in \mathcal{F}[X]$ , the sentence "Q can be splitted over  $\mathcal{F}_P$ " is true for all  $P$  belonging to  $J_Q \in \mathcal{U}$ , and so it holds in  $\prod \mathcal{F}_P/\mathcal{U}.$
- Therefore,  $\prod \mathcal{F}_P/\mathcal{U}$  is an algebraically closed field.
- Finally, we define  $\overline{\mathcal{F}}$  to be the set of elements of  $\prod \mathcal{F}_P/U$  which is a root of some  $P \in \mathcal{F}[X]$ . Clearly,  $\overline{\mathcal{F}}$  is an algebraic extension of  $\mathcal{F}$ .
- Now, suppose for the contrary that there is a polynomial in  $\overline{\mathcal{F}}[X]$  that has no root in  $\overline{\mathcal{F}}$ . Then, the root should be to expressed as a root of the polynomial of  $\mathcal F$ ("Algebraic extension" is transitive), which contradicts with the definition of  $\overline{\mathcal{F}}$ .
- Therefore,  $\overline{\mathcal{F}}$  is an algebraic closure of  $\mathcal{F}$ .

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## Theorem (Frayne-Morel-Scott)

A class of structures K is an elementary class ( $EC_{\Delta}$ ) iff it is closed under elementary equivalences and ultraproduct.

### Proof.

- ( $\Rightarrow$ ) is clear. To show ( $\Leftarrow$ ), suppose that K is closed under elementary equivalences and utraproducts. Let  $T = \{\sigma : \forall \mathfrak{A} \in \mathcal{K}, \mathfrak{A} \models \sigma\}$  and we claim  $\mathcal{K} = \text{Mod}(T)$ .  $\mathcal{K} \subseteq Mod(T)$  is clear. To show  $Mod(T) \subseteq \mathcal{K}$ , we take any  $\mathfrak{B} \in Mod(T)$ . Let I be the set of finite subsets of  $\text{Th}(\mathfrak{B})$ .
- By way of contradiction, assume there is an  $i \in I$  such that  $\forall \mathfrak{A} \in \mathcal{K}(\mathfrak{A} \not\models i)$ . Suppose  $i = {\varphi_1, \ldots, \varphi_n}$ . Since for any  $\mathfrak{A} \in \mathcal{K}$ ,  $\mathfrak{A} \models \neg \varphi_1 \vee \cdots \vee \neg \varphi_n$ , we have  $\neg \varphi_1 \vee \cdots \vee \neg \varphi_n \in T$ . Since  $\mathfrak{B} \models T$ , we have  $\mathfrak{B} \models \neg \varphi_k$  for some  $k \in i$ , which contradicts  $\varphi_k \in i \subseteq \text{Th}(\mathfrak{B})$ . Therefore, for any  $i \in I$ , there exists  $\mathfrak{A}_i \in \mathcal{K}$  such that  $\mathfrak{A}_i \models i.$
- We can construct a model  $\mathfrak{A}$  of  $T = \text{Th}(\mathfrak{B})$  by ultraproduct as in the proof of compactness theorem. Then since K is closed under ultraproducts, we have  $\mathfrak{A} \in \mathcal{K}$ . Moreover, because K is closed under elementary equivalence,  $\mathfrak{A} \equiv \mathfrak{B}$  implies  $\mathfrak{B} \in \mathcal{K}$

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### Definition

 $\prod \mathfrak{A}_i/\mathcal{U}$  is called an **ultrapower** of  $\mathfrak{A}$ , denoted by  $\mathfrak{A}^I/\mathcal{U}$ , if  $\mathfrak{A}_i = \mathfrak{A}$  for each  $i \in I$ .

Let  $\lambda i.a$  denote a function which always takes the value a. For  $a \in |\mathfrak{A}|$ , we put

$$
^*a = [\lambda i.a] \in |{\mathfrak A}^I/{\mathcal U}|
$$

and define a function  $d:|\mathfrak{A}|\to|\mathfrak{A}^I/\mathcal{U}|$  by  $d(a)=\text{*}a$ , which is called a canonical embedding.

### Definition

An embedding  $\phi : \mathfrak{A} \to \mathfrak{B}$  is said to be **elementary** if  $\phi(\mathfrak{A}) \prec \mathfrak{B}$ .

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### Theorem

Let  $\prod \mathfrak{A}^I/\mathcal{U}$  be an ultrapower of  $\mathfrak{A}$ . Then the canonical embedding  $d: |\mathfrak{A}| \to |\mathfrak{A}^I/\mathcal{U}|$  is elementary. In particular,  $\mathfrak{A}\equiv \mathfrak{A}^{I}/\mathcal{U}$ .

**Proof.** For any formula  $\varphi(x_1, \ldots, x_n)$  and  $a_1, \ldots, a_n \in |\mathfrak{A}|$ , by Los theorem,

$$
\mathfrak{A}^{I}/\mathcal{U} \models \varphi({}^{\ast}a_{1}, \ldots, {}^{\ast}a_{n}) \Leftrightarrow \{i \in I : \mathfrak{A} \models \varphi(a_{1}, \ldots, a_{n})\} \in \mathcal{U}
$$

$$
\Leftrightarrow \mathfrak{A} \models \varphi(a_{1}, \ldots, a_{n}).
$$

Thus,  $d$  is an elementary embedding. Since  $d(\mathfrak{A}) \cong \mathfrak{A}, \, \mathfrak{A} \equiv \mathfrak{A}^I/\mathcal{U}.$ 

### Theorem (Keisler-Shelah)

 $\mathfrak{A}\equiv\mathfrak{B} \Leftrightarrow$  There exist an  $I$  and a ultrafilter  $\mathcal U$  such that  $\mathfrak{A}^I/\mathcal U\cong\mathfrak{B}^I/\mathcal U.$ 

**Proof.** ( $\Leftarrow$ ) is derived from the last theorem. The proof of ( $\Rightarrow$ ) is omitted since it is too technically involved. See Model Theory: Third Edition - C.C. Chang, H. Jerome Keisler for details.

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Assuming the Keisler-Shelah theorem, we obtain the following.

### **Corollary**

The structural class K is the elementary class ( $EC_{\Delta}$ ) iff the following two conditions hold.

- (1)  $K$  is closed under ultraproducts and isomorphisms.
- $(2) \mathfrak{A}^{I}/U \in \mathcal{K} \Rightarrow \mathfrak{A} \in \mathcal{K}$ . (It is closed under inverses of ultrapower).

**Proof.** To show the sufficient condition, we prove  $\mathcal K$  is closed by elementary equivalence.

- Let  $\mathfrak{A} \equiv \mathfrak{B}$  and  $\mathfrak{A} \in \mathcal{K}$ .
- By the Keisler-Shelah theorem, there is an ultrapower  $\cal U$  such that  $\mathfrak{A}^{I}/\cal U\cong \mathfrak{B}^{I}/\cal U.$
- Since  $K$  is closed under ultraproduct,  $\mathfrak{A}^{I}/\mathcal{U} \in \mathcal{K}$ .
- Because K is closed under isomorphisms,  $\mathfrak{B}^{I}/U \in \mathcal{K}$ .
- Moreover, by condition (2), we have  $\mathfrak{B} \in \mathcal{K}$ .

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• A class K of structures in a language L is called a **projective class** or **pseudo-elementary class**, denoted  $K \in \text{PC}_{\Delta}$ , if there exists an elementary class  $\mathcal{K}'\in\mathrm{EC}_{\Delta}$  in an extended language  $\mathcal{L}'\supseteq\mathcal{L}$  such that

 $\mathcal{K} = \{ \mathfrak{A} : \mathfrak{A}$  is a reduct of a model in  $\mathcal{K}'$  to  $\mathcal{L} \}.$ 

- For example, the class of orderable groups is a projective class.
- It is easy to see that  $PC_{\Delta}$  is also closed under ultraproducts and isomorphisms. Various characterizations are also known for PC∆.
- The following one is particularly interesting, and so important as it allows us to derive Craig's interpolation theorem.

### Theorem

If K,  $K' \in PC_\Lambda$  and  $K \cap K' = \emptyset$ , then there exists  $\mathcal{J} \in EC$  such that  $K \subseteq \mathcal{J}$  and  $\mathcal{J} \cap \mathcal{K}' = \emptyset$  where  $\mathcal{J} \in EC$  means that  $\mathcal{J} = Mod(\{\sigma\})$  with a single sentence  $\sigma$ .

### Homework Problem

Show that  $K$  is finitely axiomatizable iff both  $K$  and its complement are closed under ultraproducts and elementary equivalence.

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## Non-standard analysis

- <span id="page-19-0"></span>• Using ultrapowers, we can construct a large non-standard structure that properly includes a common standard structure such as natural numbers, real numbers, and function spaces as elementary substructures.
- In particular, a non-standard model of real numbers includes infinities and infinitesimals as elements, and thus provides the first rational model for Leibniz's style of infinitesimal analysis.
- Non-standard methods have been applied to various fields of mathematics. In particular, its application to analysis is called non-standard analysis.
- From now on, we fix a non-principal ultrafilter U on the natural numbers  $\omega$  (= N) and denote the ultrapower  $\prod \mathfrak{A}^I/\mathcal{U}$  of a structure  $\mathfrak A$  by  ${}^*\mathfrak{A}.$
- As shown before, there is a natural embedding  $d(a) = *a$  from  $\mathfrak A$  to \* $\mathfrak A$ . Identifying  $\mathfrak A$ and its image  $d(\mathfrak{A})$ ,  $\mathfrak A$  can be regarded as an elementary substructure of  $^* \mathfrak A$ .
- The structures like  $\mathfrak{N} = (\mathbb{N}, +, \cdot, 0, 1, <)$  and  $\mathfrak{R} = (\mathbb{R}, +, \cdot, 0, 1, <)$  etc. are called standard models. \* $\mathfrak{N},$  \* $\mathfrak{R},$  etc. are called their non-standard models.

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- A standard model and its non-standard counterpart can not be distinguished in terms of elementary (first-order) propositions. But, there might be fundamental properties that cannot be expressed in an elementary manner, e.g., the Archimedean property.
- An ordered field  $\mathfrak A$  is **Archimedian** if for any positive elements  $a, b \in A$  there exists a sufficiently large natural number  $n \in \mathbb{N}$  such that  $b < a + a + \cdots + a$  (n times).

### Theorem

<sup>∗</sup>R is a non-Archimedean ordered field.

### Proof.

• Since  $\Re$  is an ordered field and such a property can be described in elementary way, <sup>∗</sup>R is also an ordered field.

Claim: <sup>∗</sup>R is non-Archimedean

- Let  $s = \langle 1, 2, 3, \ldots \rangle \in |\Re^{\omega}|$  and  $N = [s] \in |{}^*\Re|$ .
- Then, for any natural number  $n \in \mathbb{N}$ , we have

$$
N > \underbrace{1 + {^*}1 + \cdots + {^*}1}_{n \text{ times}},
$$

since  $\{i : s(i) > n\} \in \mathcal{U}$ .



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## **Definition**

• An element  $a$  of  $|{}^*\mathfrak{R}|$  is infinite if  $\forall b \in \mathbb{R}$   $b < |a|$ . An element that is not infinite is said to be finite. 元 a が**無限小**(infinitesimal)であるとは,∀b (> 0) ∈ R |a| < b となることを \n oloment a of <sup>|</sup>\*®}| is **infinite** if ∀ሌ∈ ₪ ሌ < lal. An oloment that is not infinite i

<sup>∗</sup>R| の元 a が**無限大**(infinite)であるとは,∀b ∈ R b < |a| と

• The element a of  $|{}^*\mathfrak{R}|$  is **infinitesimal** if  $\forall b (> 0) \in \mathbb{R}, |a| < b$ .



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$$
N = \left[\langle 1, 2, 3, \ldots \rangle\right]
$$
 is infinite,  $1/N = \left[\langle 1/1, 1/2, 1/3, \ldots \rangle\right]$  is infinitesimal.

### $\sqrt{ }$  Problem  $\sqrt{ }$

(1) Show that the set of all infinitesimals is closed under the operations  $+$  and  $\cdot$ .

 $\sim$  Example  $\sim$ 

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(2) Show that  $a$  is infinite and  $1/a$  is infinitesimal.

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• It is easy to see that  $\approx$  is an equivalence relation and also preserves the operations of  $+$  and  $\cdot$ .

### Lemma

Definition

For a finite real number  $a \in |{}^*\mathfrak{R}|$ , there exists a unique  $b \in \mathbb{R}$  such that  $a \approx b$ .

**Proof.** Set  $b = \inf\{x \in \mathbb{R} : a < x\}$ . Uniqueness is obvious.

- Such a b in the above lemma is called the **standard part** of a and is denoted by  $st(a)$ . Thus,  $a - st(a)$  is infinitesimal.
- Every finite non-standard real number  $a$  can be uniquely represented by the sum of the standard real number  $st(a)$  and an infinitesimal.

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### Lemma

If  $s = \langle a_i \rangle \in \mathbb{R}^{\omega}$  and  $\lim a_i = a$ , then  $[s] \approx {}^*a$ .

For  $a, b \in |{}^*\mathfrak{R}|$ ,  $a \approx b \Leftrightarrow a - b$  is infinitesimal.

**Proof.** For any positive number  $\varepsilon \in \mathbb{R}$ ,  $\{i : |a_i - a| < \varepsilon\} \in \mathcal{U}$ . Therefore,  $[s] - *a$  is infinitesimal.

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### Definition

For 
$$
f : \mathbb{R} \to \mathbb{R}
$$
, we define  ${}^*f : |{}^*\mathfrak{R}| \to |{}^*\mathfrak{R}|$  as follows: for  $s \in |\mathfrak{R}^{\omega}|$ ,  
 ${}^*f([s]) = [\lambda i.f(s(i))]$ .

The well-definedness of  $*f$  follows from

$$
||s = s'|| \in \mathcal{U} \Rightarrow ||\lambda i.f(s(i)) = \lambda i.f(s'(i))|| \in \mathcal{U}.
$$

Also, \* f can be obtained from the ultrapower \* $\Re \cup \{ * f \}$  of

$$
\mathfrak{R}\cup\{f\}=(\mathbb{R},f,+,\,\bullet\,,0,1,<).
$$

### Theorem

$$
f: \mathbb{R} \to \mathbb{R} \text{ is continuous at } a \in \mathbb{R} \Leftrightarrow \text{ for any } x \approx a, \;^* f(x) \approx f(a).
$$

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• Let  $f : \mathbb{R} \to \mathbb{R}$  is continuous at  $a \in \mathbb{R}$ , and  $x = |\langle x_i \rangle| \approx a$ . Take any positive number  $\varepsilon \in \mathbb{R}$ . By the continuity of f, there exists a positive number  $\delta \in \mathbb{R}$  such that

$$
\forall y \in \mathbb{R}(|y - a| < \delta \to |f(y) - f(a)| < \varepsilon).
$$

- Therefore,  $\{i : |x_i a| < \delta\} \subseteq \{i : |f(x_i) f(a)| < \varepsilon\}.$
- Since  $x \approx a$ , we have  $\{i : |x_i a| < \delta\} \in \mathcal{U}$ .
- Hence,  $\{i : |f(x_i) f(a)| < \varepsilon\} \in \mathcal{U}$ . That is,  $*f(x) \approx f(a)$ .
- $(\Leftarrow)$

Proof.  $(\Rightarrow)$ 

- Suppose that f is not continuous at  $a \in \mathbb{R}$ .
- That is, there exists a positive number  $\varepsilon \in \mathbb{R}$  such that for any  $i \in \omega$ , there exists  $x_i$ such that

$$
|x_i - a| < \frac{1}{i+1} \land |f(x_i) - f(a)| \ge \varepsilon
$$

• Let  $x = \langle x_i \rangle$ . Then  $x \approx a$ ,  $|f(x) - f(a)| \ge \varepsilon$ . In other words,  $f(x) \not\approx f(a)$ .  $\Box$ 

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• Let  $f : \mathbb{R} \to \mathbb{R}$  be a continuous function. By the theorem, for any finite  $a \in |{}^* \mathfrak{R}|$ ,

$$
st(*f(a)) = f(st(a)).
$$

- The relationship  $S \subseteq \mathbb{R}^n$  of  $\mathfrak R$  can be naturally extended to the relation  $^*S$  of  $^*\mathfrak R$ . In particular, <sup>∗</sup>N and <sup>∗</sup>Q can be viewed as subsets of | <sup>∗</sup>R|. Moreover, notice that  $(\mathfrak{R}, \mathbb{N}, \mathbb{Q})$  is an elementary substructure of  $({}^*\mathfrak{R}, {}^*\mathbb{N}, {}^*\mathbb{Q}).$
- Let  $N = \{(1, 2, 3, ...) \} \in {^*}\mathbb{N}$ . We consider an  $N$ -partition of  $[0, 1]$  in  ${}^* \mathfrak{R}$  as  $\{0, 1/N, \ldots, (N-1)/N, N/N\}.$
- Given a standard real number  $a$  of  $[0,1]$ , take  $i \in {}^*\mathbb{N}$  with  $i/N \leq a \leq (i+1)/N$ , and then we have  $a = st(i/N)$ . In other words, any standard real number can be expressed as a non-standard fraction.
- Based on the above observations, many theorems in analysis can be proven by using the non-standard method. Here we will give two examples.

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Theorem

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A continuous function  $f : [0, 1] \rightarrow \mathbb{R}$  has the maximum value.

**Proof.** In  $*$ R, consider the following  $*$  finite set

 $\{^*f(0),^*f(1/N),\ldots,^*f((N-1)/N),^*f(N/N)\}.$ 

If its maximum value is \*  $f(i/N)$ , f has the maximum value st(\* $f(i/N)$ ) at  $x = st(i/N)$ .

 $\sf Remark.\; Since\: ({\mathfrak R},{\mathbb N})$  is an elementary substructure of  $({}^*{\mathfrak R},{}^*{\mathbb N}),$  one can use mathematical induction on  $*N$ . For instance, it is provable that any  $*$  finite set has the maximal element.

 $\Box$ 

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### <span id="page-28-0"></span>Theorem (Peano)

Let  $f : [0, 1]^2 \to \mathbb{R}$  be a continuous function. The following differential equation has a solution

$$
dy/dx = f(x, y), y(0) = 0.
$$

**Idea of the proof**<sup>1</sup> We define  $Y : \{0, 1/N, \ldots, N/N\} \rightarrow \mathbb{R}$  inductively as follows:  $Y(k/N) =$  $\sum_{i=1}^{k-1} \varepsilon^* f(Y(i/N), i/N) \cdot 1/N.$ 

Then,  $y : [0, 1] \to \mathbb{R}$  is defined as follows: given a standard real number  $a$  of  $[0, 1]$ , take  $k \in {}^*\mathbb{N}$  with  $k/N \le a \le (k+1)/N$  and set  $y(a) = \text{st}(Y(k/N)).$ П

# Thank you for your attention!

 $1$ V. Benci and M. Di Nasso, How to Measure the Infinite: Mathematics with Infinite and Infinitesimal Numbers, World Scientific Publishing, 2019.